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METHOD AND CIRCUIT ARRANGEMENT FOR IMPROVED DATA TRANSMISSION

Field of the Invention

present

The invention relates to a method and a circuit arrangement for improved data transmission according to the preambles of Claim 1 and Claim 9, respectively.

Related Technology

In telecommunications engineering, transmission methods are known, and also used in practice, which utilize orthogonal basis functions to represent the signal to be transmitted. Such transmission methods are described, for example, in the book of R.E. Blahut, *Digital Transmission of Information*, Addison-Wesley, Reading, 1990, chapters 2 and 3.

In this case, a message signal $s(t)$ in the baseband is represented as the sum of orthogonal basis functions. In order to integrate the message $m = (m_0, m_1, m_2 \dots m_{k-1})$ - where the m_i are selected from an appropriately chosen alphabet - into the signal $s(t)$, the signal is formed as follows:

$$s(t) = m_0 f_0(t) + m_1 f_1(t) + \dots + m_{K-1} f_{K-1}(t).$$

Consequently, a message signal can be regarded as a point in a K-dimensional space and specifically is characterized by the value-tuple $(m_0, m_1, \dots, m_{K-1})$. The entirety of all permissible signals is referred to as a signal constellation. Especially popular in practice are two-dimensional signal constellations, such as the so-called 16-QAM signal constellation shown in Fig. 1 of the present application. This 16-QAM signal constellation is described, for example, in the aforementioned book on page 63. In all the signal constellations considered here, it is assumed that the minimum distance between two signal points is normalized to 1. However, the known transmission methods for efficient use of multi-level modulation methods do not yet permit the

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optimal utilization of the signal energy of signal constellations. First of all, signal constellations which are very efficient may, however, have the disadvantage that the number of signal points is not a power of two and, secondly, frequently employed signal constellations, such as 16-QAM, cannot yet be used in simple and optimal manner to transmit low data rates.

Fundamental theoretical investigations on improved data transmission and on the efficient use of multi-level modulation methods employing orthogonal basis functions to represent a signal to be transmitted and using, for example, the known Huffman method as a source coding method have been published in F.R. Kschischang, S. Pasupathy, "Optimal Nonuniform Signaling for Gaussian Channels", IEEE Transactions on Information Theory, Vol. 39, No. 3, May 1993, pp. 281-300. However, practical implementations of these investigations in the form of circuit arrangements and/or corresponding methods for operating such circuit arrangements have not been specified.

SUMMARY of the Invention

Therefore, the object of the invention is to provide a method and a circuit arrangement for improved data transmission with efficient use of multi-level modulation methods which permit optimal use of the signal energy of signal constellations and by which frequently employed signal constellations, such as 16-QAM, can be used in simple and optimal manner to transmit lower data rates.

~~The objective of the invention with regard to the method is achieved by the method described in the characterizing part of Claim 1.~~

Further design approaches or refinements of the method are described in Claims 2 through 8 and 12 through 15.

~~The objective of the invention with regard to the circuit arrangement is achieved by the circuit arrangement described in the characterizing part of Claim 9 and a further~~

A refinement is described in the characterizing parts of Claims 10 through 12

A Further design approaches and refinements of the invention are indicated in the following detailed description

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The method and the circuit arrangement described here permit optimal utilization of the signal energy of signal constellations. This can be of advantage for technical applications in two respects. First of all, signal constellations which are very efficient, but have the disadvantage that the number of signal points is not a power of two, can now be adapted in a simple manner to data formats used in practice, such as a bit sequence. Secondly, frequently employed signal constellations, such as 16-QAM, can be used in simple and optimal manner to transmit lower data rates. Thus, the 16-QAM signal constellation can be used to transmit on average 3 bits per signal point, instead of the usual 4 bits per signal point. This can be technically useful, for example, in order to switch to 16-QAM with optimized probabilities in existing transmitters and receivers using, for instance, 8-QAM with the points { (+-1/2, +-1/2), (+- (1+√3)/2, 0), (0, +- (1+√3)/2) as signal constellation (i.e. 3 bits per signal point), accompanied by simultaneous power gain of approx. 1 dB.

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W³ ~~The method has a further characteristic which can be put to advantageous use. This is the recoding, particularly simple to accomplish, if the input data stream is a uniformly distributed sequence, especially a bit sequence. The recoding can then be accomplished using a loss-free decompression method such as the Huffman method. Accordingly, the inverse recoding operation on the receiver end is carried out using the corresponding compression method. A uniformly distributed sequence or bit sequence is obtained, for example, by encryption. This means that the possibly bothersome guarantee or generation of such a sequence can be achieved by the addition of a value-added operation, namely encryption. Since encryption will play an ever greater role in future transmission systems and is already being supplied today along with many systems, the new method is particularly practical. In the recoding of~~

the incoming source bit stream to the signal points which are transmitted through the channel, use is made of a temporary register as buffer that serves to adapt the bit rate, which fluctuates as a function of time due to the transmission through the channel, to the bit rate of the source data. In an implementation of the circuit, this temporary register has a defined fixed length. In practice, this leads to the problem of a "buffer overflow". To solve this problem, the suggestion here is to select the channel data rate to be greater than the source data rate, it being advantageous to select the channel data rate to be slightly greater than the source data rate. In this manner, it is possible, with relatively little effort and expense, to specify the length of the temporary register or the length of the buffer so that there is only a negligibly small (known) probability of an overflow. When working with a channel data rate which is greater than the source data rate, it may occur that the channel will be ready to transmit information which is not yet available from the source. This effect is utilized here, for example, by transmitting synchronization data instead of the source data. A further solution involves transmitting other header or user data as well, instead of the synchronization data. The greater the channel data rate, the shorter it is possible to select the temporary register.

Further advantages, features and application possibilities of the present invention are revealed from the following exemplary embodiments, which are described with reference to the drawing and tables.

~~The terms and reference numerals used in the appended list of reference numerals are used in the specification, in the claims, in the abstract and in the drawing.~~

Present Brief Description of the Drawings
Hereinbelow, the invention is described in greater detail on the basis of exemplary embodiments with reference to the drawing, in which:

a graphical representation of
Fig. 1 shows a 16-QAM signal constellation;

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Fig. 2 ^{a graphical representation of} shows a hexagonal signal constellation;

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Fig. 3+6 show a block diagram of a circuit arrangement ~~used~~ for improved data transmission with the aid of the efficient use of multilevel modulation methods;

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Fig. 4 shows a table 1, which indicates the probabilities p_1, p_2, p_3, p_4 for the signal ~~shown~~ points in Fig. 2; ~~and~~

10 Fig. 5 shows a table 2, which represents the mapping of the binary data to the signal points and vice versa; ~~and~~

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As ^{stated above} ~~was already stated~~, known transmission methods employ orthogonal basis functions to represent the signal to be transmitted. In this case, a message signal $s(t)$ is represented as the sum of orthogonal basis functions. A message signal can be regarded as a point in a K-dimensional space. The entirety of all permissible signal points is referred to as a signal constellation, the "16-QAM" signal constellation shown in Fig. 1, which represents one of the two-dimensional signal constellations, being especially popular.

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If a signal constellation has a total of M signal points, M_j of each of which have the signal energy E_j , and if the probability for the occurrence of such a signal point is equal to P_j , then, by setting the probabilities according to the formula given below, one obtains the values which are optimal according to power/information rate for that power. The value L indicates how many different energy levels occur in total.

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$$p_j = p_1 (p_L/p_1)^{(E_j-E_1)/(E_L-E_1)} \quad j = 1, 2, \dots, L \text{ and } E_{j+1} > E_j$$

Given here as an example is the hexagonal signal constellation in Fig. 2. For reasons of normalization, the minimum distance between the signal points is selected as one. Here, there are $L = 4$ energy levels. $E_1 = 0$, $E_2 = 1$, $E_3 = 3$ and $E_4 = 4$. There is one signal point with signal energy zero ($M_1 = 1$) and 6 signal points each with signal energies 1, 3 and 4, i.e., $M_2 = M_3 = M_4 = 6$.

For example, to map a data stream with a defined probability distribution to the corresponding signal points, use is made of a loss-free data compression algorithm, such as the Huffman method. This data compression algorithm ensures that the corresponding signal points occur with the aforementioned probability. The Huffman method is described, for example, in D.A. Huffman, "*A Method for the Construction of Minimum Redundancy Codes*", Proc. IRE, Vol. 40, Sept. 1952, pages 1098-1101. In the following example, a binary bit sequence, in which the probability for ones and zeros is identical and in which the bits are statistically independent, is recoded, and specifically in such a way that with the signal constellation shown in Fig. 2 with 19 signal points, on average $H = 4$ bits per signal point can be efficiently transmitted. From the indicated table 1 according to Fig. 4, one then obtains the probabilities for the occurrence of the individual signal points. The use of a data compression method leads to a correspondence of the kind found, for example, in table 2. With the correspondence shown in table 2 according to Fig. 5, one comes very close to the optimal mean signal energy $E_s = 1.7224$. One obtains a mean signal energy of 1.8125. By comparison, the conventional 16-QAM signal constellation has a mean signal energy of 2.5. In other words, in comparison with the known 16-QAM, one obtains an improvement of $10 \lg (2.5/1.8125)$ dB, that is, approximately 1.4 dB, with this simple method. With more complex correspondences, one can come as close as one wishes to the optimal value. For the purpose of illustration, with the above correspondence, the bit sequence 01110100001111100111011110001 produced with a coin would then be transmitted with the signal points $Z_{32} Z_1 Z_{25} Z_{23} Z_{25} Z_{21} Z_{24} Z_{25}$. Z_1 is the signal point with energy zero; Z_{2j} , where $j = 1..6$, are the signal points with energy 1; Z_{3j} are the signal points with energy 3 and Z_{4j} are the points with energy 4.

The decoding after transmission follows accordingly.

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5 In the following, ~~the construction of a circuit arrangement for implementing the~~
~~above-described method~~ is explained in greater detail with reference to ~~a~~ ^{the} block
diagram according to Fig. 3.

10 It is assumed that a data source 1 supplies a data stream 2. A recoder 3 then ensures
that a modulator 4 selects the corresponding signal points with the correct probability.
After data stream 2 has been transmitted via a transmission channel 5, there follows,
after a downstream demodulator 6, the corresponding inverse operation with the aid of
10 ^{inverse} a recoder 7, whereupon data stream 2 finally arrives at a data sink 8. The respective
data stream 2 is depicted on the connecting/transmission lines or channels between
components 1, 3 through 8 by arrow points on the respective lines/channels.

15 In the recoding of the incoming source bit stream to the signal points which are
transmitted through the channel, a temporary register (not shown) is inserted as buffer
that is used to adapt the bit rate, which fluctuates as a function of time due to the
transmission through the channel, to the bit rate of the source data. In each
implementation in the form of a circuit, the temporary register or buffer has a defined
20 length, this possibly leading in practice to a problem of overflow. This problem can
be solved by selecting the channel data rate to be somewhat greater than the source
data rate. It is therefore possible, with relatively little effort and expense, for the
buffer length or the temporary-register length to be specified so that there is only a
negligible small, ~~(known)~~ probability of a buffer overflow or temporary-register
25 overflow.

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If the channel data rate is selected to be slightly greater than the source data rate, it
may occur in the practical implementation of a circuit arrangement that the channel
will be ready to transmit information which is not yet available from the source.

This effect can be advantageously utilized by transmitting synchronization data instead of the source data. Furthermore, other header or user data can also be transmitted instead of the synchronization data. The greater the channel data rate in relation to the source data rate, the shorter it is possible to select the temporary register or temporary buffer.

A further solution to this problem of buffer overflow or underflow is to provide two or possibly more recoding tables in recoder 3, the one table leading to a channel data rate which is greater than the source data rate and the other table leading to a channel data rate which is lower than the source data rate. Recoder 3 can then be controlled as a function of the state of the temporary storage. That is, if the temporary storage is in danger of overflowing, the channel data rate is selected which is greater than the source data rate. In the opposite case, if there are almost no data left in the temporary storage, the channel data rate is selected which is lower than the source data rate.

Fig. 6, which is derived from Fig. 3, shows the most common case, which includes the above possibility. The possibility of controlling recoder 3 as a function of temporary storage 9 is indicated in Fig. 6 by the broken line from the temporary storage with control unit/processing unit 9 to recoder 3. Likewise shown in Fig. 6 is an optional second data source 1' (for the special case when the rate of this data source is equal to zero, this source disappears). As described above, second data source 1' permits the transmission of additional data. The broken lines from recoder 3 to second data source 1' indicate the manner in which, for example, check data can be integrated into the method for the purpose of correcting errors. The source data rate and the rate of the generated error detection characters together must, on average, not exceed the mean channel data rate. Analogous to second data source 1' and the temporary storage with control/processing unit 9, a second data sink 8' and a temporary storage with control/processing unit 9' are inserted between data sink 8 and inverse recoder 7.

To further improve the method, it is possible to use special coding methods which

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may be used
have been designed, for example, for QAM or hexagonal signal constellations and
~~which can be found in the articles~~ K. Huber, "*Codes over Gaussian Integers*", IEEE
Transactions on Information Theory, Vol. 40, No. 1, January 1994, pp. 207-216 and
K. Huber, "*Codes over Eisenstein-Jacobi Integers*", Finite Fields: Theory,
5 Applications and Algorithms, (Las Vegas 1993), Contemporary Math. Vol. 168,
American Math. Society, Providence, RI, pp. 165-179 as well as K. Huber, "*Codes
over Tori*", IEEE Transactions on Information Theory, Vol. 43, No. 2, March 1997,
pp. 740-744. ✓^{12/13}

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List of reference numerals

| | |
|------|------------------------------------------------|
| 1,1' | Data source |
| 2 | Data stream |
| 3 | Recoder |
| 4 | Modulator |
| 5 | Channel |
| 6 | Demodulator |
| 7 | Inverse recoder |
| 8,8' | Data sink |
| 9,9' | Temporary storage with control/processing unit |